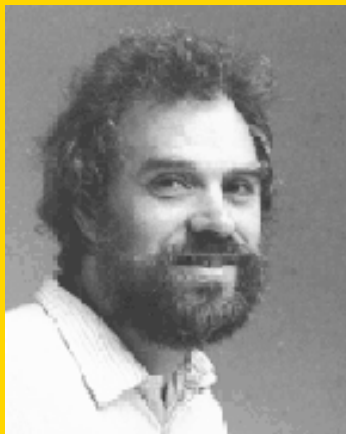


Abstract

Jacobi-Davidson QZ is an efficient and robust algorithm for finding one or a few eigenvalues with specific properties (largest, with largest real part, closest to zero) and for finding the associated eigenvalues. The method applies to all sorts of square matrices: symmetric, non-symmetric, non-normal, complex, etc.. In comparison with other methods it is specifically powerful for large sparse matrices. There are extensions for solving quadric eigenproblems and for even higher order polynomial eigenproblems.



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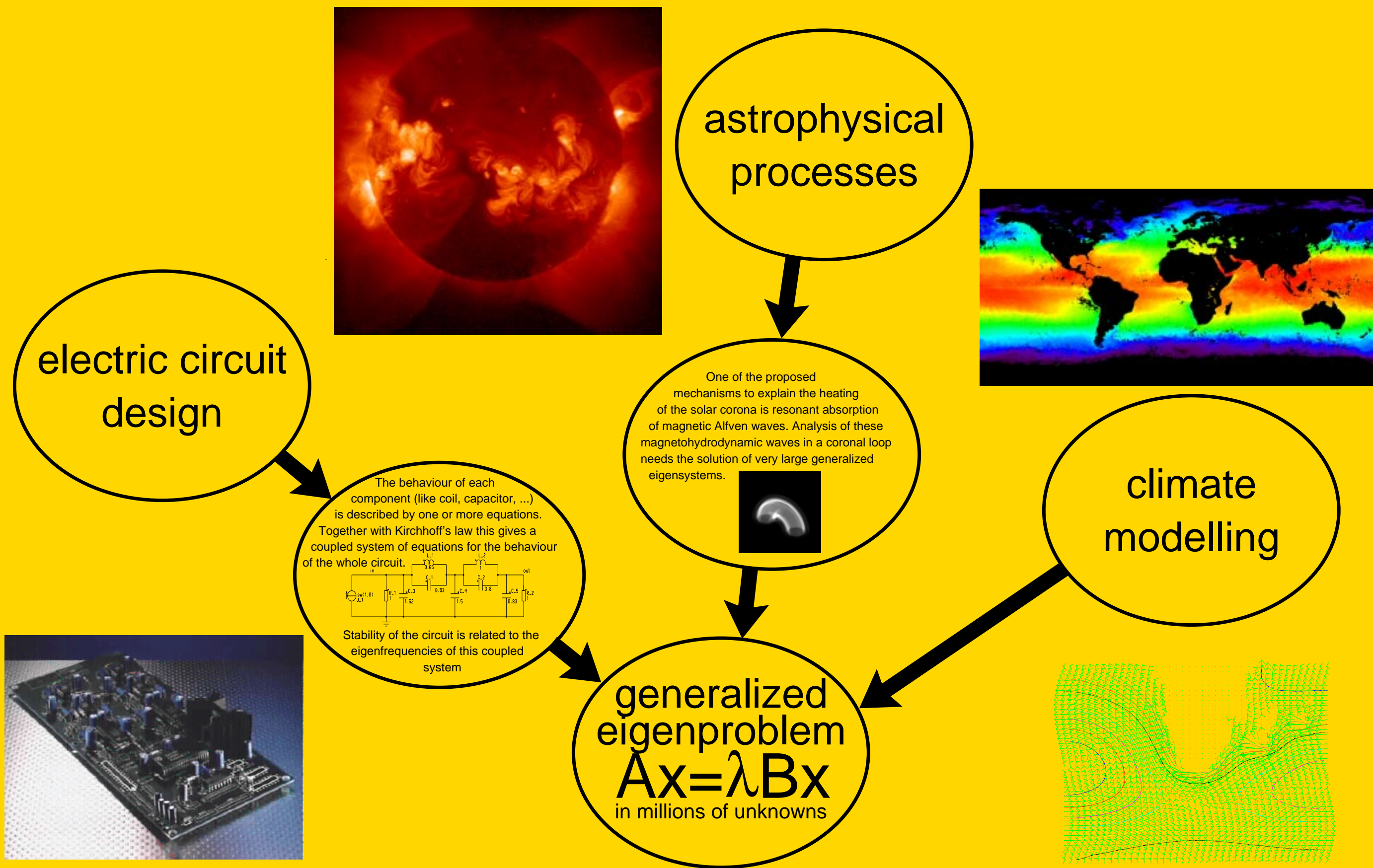
Henk van der Vorst



Menno Genseberger

Jacobi-Davidson QZ

an efficient tool for solving Generalized Eigenproblems



For given square n by n matrices A and B and some value $\tau \in \mathbb{C}$, find scalars $\lambda \in \mathbb{C}$ close to the target value τ and vectors x for which $Ax = \lambda Bx$ λ is an eigenvalue, x is an eigenvector associated with λ .

High resolution discretizations of the modelling partial differential equations lead to generalized eigenproblems of very high dimension. In the project, there was a need for solvers of eigenproblems of dimension 600 000 and more. The conventional approaches could not cope well with such high dimensions.

Details on Jacobi-Davidson

Jacobi-Davidson searches for an approximate eigenvector in a low dimensional **search subspace** V_k . It expands this subspace with one vector if the obtained approximation is not accurate enough. This procedure is repeated until convergence.

The low dimensionality allows efficient computations.

Fundamental questions

- How to **extract** approximate eigenvectors from V_k ? *Answer: Galerkin*
- How to **expand** V_k ? *Answer: Newton*

Both answers are optimal with respect to the information that is available. Jacobi-Davidson is (inexact) Newton with subspace acceleration in which all steps are optimally tailored for (generalized) eigenvalue problems.

EXTRACT Petrov-Galerkin

Find $u \in V_k = \{v_1, \dots, v_k\}$ such that

$$r = (A - \theta B)u \perp \mathcal{W}_k = \{w_1, \dots, w_k\}.$$

(θ, u) is the **Petrov pair** w.r.t. the **search** subspace V_k and the **test** subsp. \mathcal{W}_k .

In practice. $k \ll n$
With both $V_k \equiv [v_1 | v_2 | \dots | v_k]$ and $\mathcal{W}_k \equiv [w_1 | w_2 | \dots | w_k]$ orthonormal, and $u = V_k y$, we have a low dimensional generalized eigenproblem

$$(W_k^* A V_k - \theta W_k^* B V_k) y = 0.$$

Apply QZ for the low dimensional reduced matrices $W_k^* A V_k$ and $W_k^* B V_k$.

EXPAND Inexact Newton

Compute the **residual** vector r and an auxiliary vector $z = \bar{z} / \|z\|_2$ where

$$r \equiv Au - \theta Bu \quad \text{and} \quad \bar{z} \equiv \beta Au + \alpha Bu.$$

Find an approximate solution $\tilde{t} \perp u$ of the Jacobi-Davidson correction equation

$$u^* t = 0 \quad \text{and} \quad (I - z z^*) (A - \theta B) (I - u u^*) t = -r$$

Expand V_k with \tilde{t} ; orthonormalize \tilde{t} w.r.t. V_k to v_{k+1} .

Expand \mathcal{W}_k with $\tilde{s} \equiv \alpha A \tilde{t} + \beta B \tilde{t}$; orthonormalize \tilde{s} w.r.t. \mathcal{W}_k to w_{k+1} .

The Jacobi-Davidson corr. eq. is a Jacobian system of a Newton process. Therefore: quadratic convergence when solved exactly; but, due to the subspace acceleration, also fast convergence with inexact solutions.

In practice. Any linear solver can be used to solve the correction equation. Complication: the correction equation is singular in \mathbb{C} . However: the correction equation is non-singular from u^\perp to z^\perp . Preconditioned iterative Krylov subspace solvers, as preconditioned GMRES, are very suitable for such situations: all iterates are orthogonal to u if that is the case for the initial guess. Often already a few steps lead to fast convergence.

Preconditioners

If M is a preconditioner for $A - \theta B$, then M_p is a preconditioner for A_p , where

$$M_p \equiv (I - z z^*) M (I - u u^*) \quad \text{and} \quad A_p \equiv (I - z z^*) (A - \theta B) (I - u u^*).$$

The action of the preconditioned operator $M_p^{-1} A_p$ on vectors s can be efficiently incorporated in solvers as GMRES:

First solve y from $My = z$ and compute $\mu \equiv u^* y$. *Only once per linear solve!*
Then, in each GMRES step, compute $t \equiv M_p^{-1} A_p s$ as

- $s \leftarrow (A - \theta B)s$
- Solve $M s' = s$
- $t \leftarrow s' - y(\mu^{-1}(u^* s'))$

Only one (cheap) projection!

Note that M can also be used for other nearby θ . It pays to make a good preconditioner.

Jacobi-Davidson QZ

Restart. Increasing storage or computational overhead, for increasing dimension k of the subspaces, may make it necessary to restart (if, say, $k > 30$). With a restart by a single vector, valuable information that is contained in the remaining part of the subspace may be discarded, leading to a slowdown of speed of convergence. Therefore, Jacobi-Davidson is restarted with a space spanned by a small number (say, 5) of Petrov vectors with Petrov values closest to the target τ . These are the vectors that contain the most valuable information for the wanted eigenpair.

Deflation. When a Petrov value is close enough to an eigenvalue, the remaining part of the current subspace will already have rich components in nearby eigenpairs. This information is used as basis for a subspace for the computation of the next eigenvector. In order to avoid that old eigenvectors will re-enter the computational process, the new search vectors are explicitly orthogonalized to the computed eigenvectors. This technique is called *deflation*.

Partial QZ-decomposition. For stability and for easy computation, working with orthonormal basis is preferred. Instead of eigenvectors, a partial QZ-decomposition is computed:

$$A Q_m = Z_m S_m \quad \text{and} \quad B Q_m = Z_m T_m.$$

Here the columns of the n by m matrix Q_m form an orthonormal basis of the space spanned by m wanted eigenvectors. The matrix Z_m is also n by m and orthonormal. S_m and T_m are square m by m upper triangular. As for the full QZ-decomposition, eigenvalues with associated eigenvector can easily be extracted from the partial decomposition.

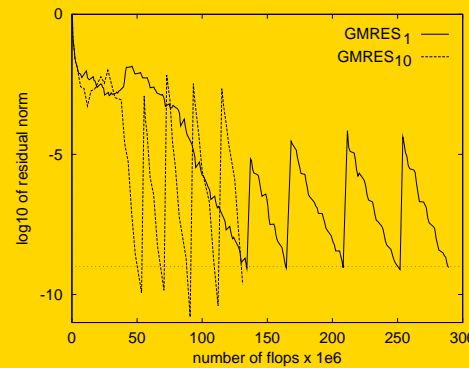
Jacobi-Davidson QZ. The use of the QZ-decomposition for the reduced matrices allows easy accommodation of restarts and deflation and leads in a natural way to a partial QZ-decomposition. The resulting scheme is called **Jacobi-Davidson QZ** and can be viewed as a truncated form of the QZ-algorithm for large generalized eigenvalue problems.

Solving the correction equation accurately: two small examples

The examples here are of relative low order and serve as illustration of the effect of solving of the correction equation more accurately. In both examples the five rightmost eigenvalues are computed. The figures show the \log_{10} of the norm of the residual as a function of the computational costs (expressed in mega floating point operations). Every time when the norm is less than 10^{-8} , an eigenvalue is determined and the iteration is continued with deflation for the next eigenvalue. The maximum dimension of the search subspace is 15. After restart the dimension is 10. One and ten steps of GMRES were used to compute an approximate solution of the correction equation.

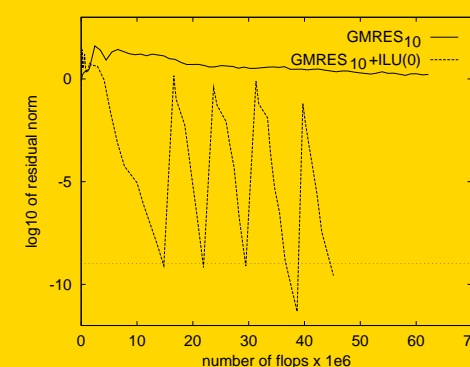
Both examples show that it pays to solve the correction equation more accurately. Although the projections become more expensive after convergence of a few eigenvectors, the computational costs to find the next eigenvectors does not increase. Apparently, the speed of convergence improves.

Example 1. The eigenproblem of order 4096 comes from the stability analysis of an integrated circuit. No preconditioner has been used.



Note that after an initial phase with two small bumps, JDQZ converges quite fast. For the next two eigenvalues there is no such initial stagnation. Apparently, in the iteration for the first eigenvalue, components for the next eigenvector are already collected in the search subspace.

Example 2. The eigenproblem ("Brus-selator") of order 2000 models the stability of concentration waves for reaction and transport interaction of chemical solutions in a tubular reactor.

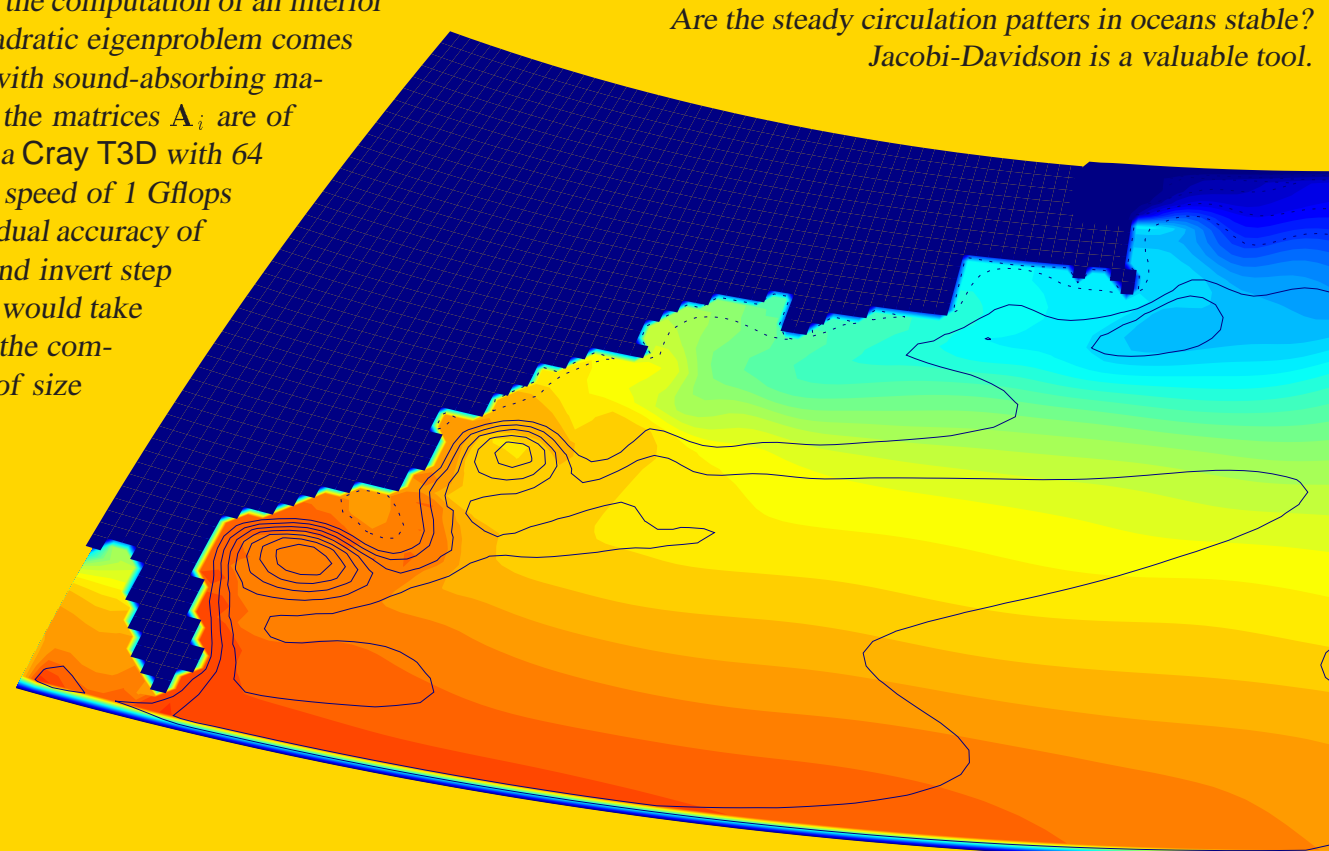


The preconditioner ILU(0) (incomplete LU of $A - B$) for GMRES has been constructed only once, and appears to remain efficient also for eigenvalues further away from the target $\tau = 1$.

A large example

Jacobi-Davidson can also be used for interior eigenvalues and for higher order polynomial eigenproblems. As an illustration, we discuss the computation of an interior eigenvalue of $A_0 x + \lambda A_1 x + \lambda^2 A_2 x = 0$. The quadratic eigenproblem comes from modelling sound waves in a room with a wall with sound-absorbing material. For the selected finite element discretization the matrices A_i are of size 274 625. Diagonal preconditioner was used. On a Cray T3D with 64 processors, it took 93.4 seconds with computational speed of 1 Gflops to compute the wanted eigenvalue with a relative residual accuracy of 10^{-6} . At peak performance (9.6 Gflops), one shift and invert step would take more than 10 minutes, the QZ-algorithm would take more than a month (if it would fit in the memory of the computer). The algorithm scales well as, for matrices of size 136 161, is shown by the table.

Number of processors	wall clock time
16	167.7
32	82.5
64	41.5



Are the steady circulation patterns in oceans stable? Jacobi-Davidson is a valuable tool.

Conventional Approaches

I. The QZ-algorithm

- The **QZ-algorithm** produces square orthogonal matrices Q and Z for which

$$Z^* A Q = S \quad \text{and} \quad Z^* B Q = T$$

are upper triangular. QZ codes are available in software packages as LAPACK.

- Then $Ax = \lambda Bx$ for $\lambda = S_{i,i}/T_{i,i}$ and $x = Qy$ where $y \neq 0$ satisfies the upper triangular system $(S - \lambda T)y = 0$.

Only feasible for low dimensional problems ($n \lesssim 1000$)

II. Shift and Invert Arnoldi

- Shift and invert** to a standard eigenvalue problem:

$$[(A - \tau B)^{-1} B] \bar{x} = \bar{\lambda} \bar{x}$$

Then $Ax = \lambda Bx$ for $\lambda = \tau + 1/\bar{\lambda}$ and $x = \bar{x}$.

- Arnoldi** (or Lanczos) for the shifted and inverted problem can be used to compute the eigenvalue(s) λ closest to τ .

Requires exact solutions of systems $(A - \tau B)t = s$. This is costly or unfeasible for high dimensions

New Approach Jacobi-Davidson

Characteristics

- Applicable for high dimensional problems. The matrices may be complex and non-symmetric.
- No exact solutions are needed of systems involving A and B .
- Iterative method that uses the following high dimensional operations

- vector updates
- DOT products
- multiplications by A and B (no inversion or exact solves)
- application of some preconditioner

plus the QZ-algorithm on low dimensional problems.

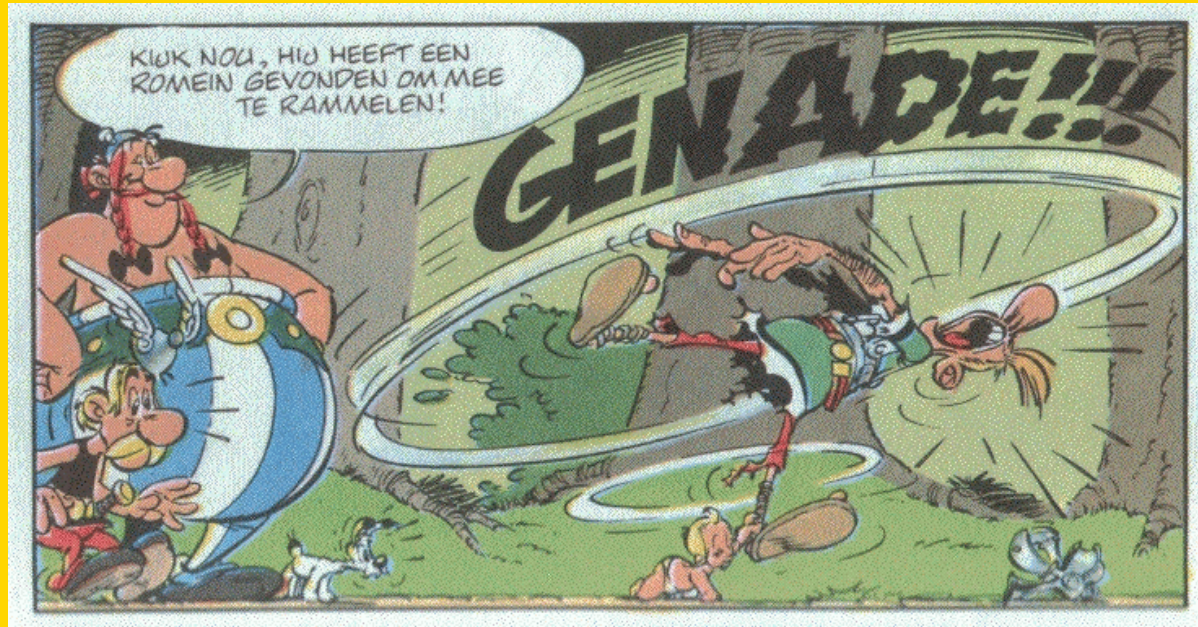
High dimensional operations are well parallelizable (how well depends on the sparsity structure of the matrices A and B and on the selected preconditioner).

- Fast convergence.
- Information from simpler models (preconditioners) can be exploited to speed up convergence.
- Can also find eigenvalues with associated eigenvectors in the interior of the spectrum (= set of eigenvalues).
- There are extensions (JDQZ) for computing a specific portion of the spectrum.
- There are extensions for quadratic and higher order polynomial eigenproblems.
- Convergence to specific eigenvalues can be guided.

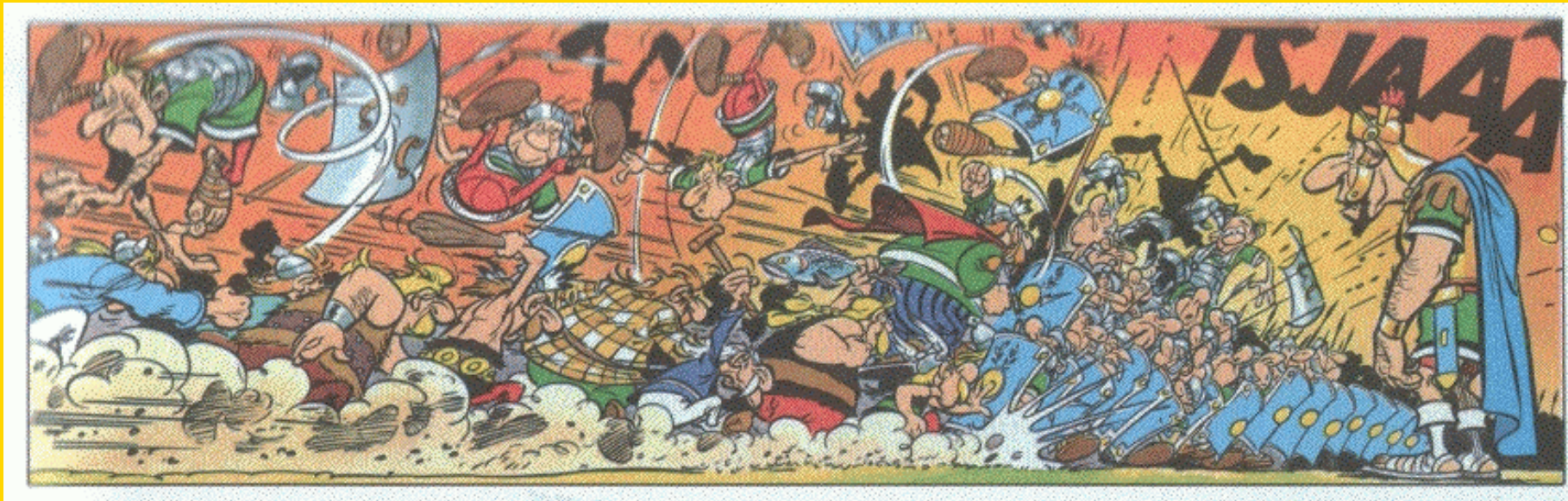
A new and robust tool for finding eigenfrequencies



... reduced models can be exploited ...



... operations are well parallelizable ...



cartoons taken from Asterix & Obelix

More Details

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